

# **Tutorial: Extending Loop Transformation Frameworks to Irregular Applications (aka Math for Irregular Codes)**

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# Transformation Frameworks

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## Motivation

- Intermediate representations for computations
- Transformation specifications and code generation after transformation
- Composition of transformations
- Data dependence representation and thus legality checks for composed transformations

## Covering in Today's Tutorial

- Polyhedral Model for representing computations with affine loop bounds and array accesses
- Polyhedral compilation for sparse-immutable computations
- Sparse Polyhedral Framework (SPF) for sparse matrix/tensor computations with indirect array accesses
- PolyRec Framework for recursive irregular computations

# Polyhedral model: Loop transformations improve performance!

## Example (dgemm)

```
/* C := alpha*A*B + beta*C */
for (i = 0; i < ni; i++)
  for (j = 0; j < nj; j++)
S1:   C[i][j] *= beta;
for (i = 0; i < ni; i++)
  for (j = 0; j < nj; j++)
    for (k = 0; k < nk; ++k)
S2:   C[i][j] += alpha * A[i][k] * B[k][j];
```

- ▶ Loop transformation: *permute(i,k,S2)*

**Execution time (in s) on this laptop, GCC 4.2, ni=nj=nk=512**

version	-O0	-O1	-O2	-O3 -vec
original	1.81	0.78	0.78	0.78
permute	1.52	0.35	0.35	0.20

<http://gcc.gnu.org/onlinedocs/gcc-4.2.1/gcc/Optimize-Options.html>

# Another Example: FDTD

## Example (fdtd-2d)

```
for(t = 0; t < tmax; t++) {
  for (j = 0; j < ny; j++)
    ey[0][j] = _edge_[t];
  for (i = 1; i < nx; i++)
    for (j = 0; j < ny; j++)
      ey[i][j] = ey[i][j] - 0.5*(hz[i][j]-hz[i-1][j]);
  for (i = 0; i < nx; i++)
    for (j = 1; j < ny; j++)
      ex[i][j] = ex[i][j] - 0.5*(hz[i][j]-hz[i][j-1]);
  for (i = 0; i < nx - 1; i++)
    for (j = 0; j < ny - 1; j++)
      hz[i][j] = hz[i][j] - 0.7* (ex[i][j+1] - ex[i][j] +
        ey[i+1][j]-ey[i][j]);
}
```

- ▶ Loop transformation: *polyhedralOpt(fdtd-2d)*

**Execution time (in s) on this laptop, GCC 4.2, 64x1024x1024**

version	-O0	-O1	-O2	-O3 -vec
original	2.59	1.62	1.54	1.54
polyhedralOpt	2.05	0.41	0.41	0.41

# Doing such transformations by hand is NOT FEASIBLE!

## Example (fddd-2d tiled)

```
for (c0 = 0; c0 <= (((ny + 2 * tmax + -3) * 32 < 0?((32 < 0?-((-ny + 2 * tmax + -3) + 32 + 1) / 32) : -((-ny + 2 *
tmax + -3) + 32 - 1) / 32))) : (ny + 2 * tmax + -3) / 32)); ++c0) {
    #pragma omp parallel for private(c3, c4, c2, c5)
    for (c1 = (((c0 * 2 < 0?-(-c0 / 2) : ((2 < 0?(-c0 + -2 - 1) / -2 : (c0 + 2 - 1) / 2)))) > (((32 * c0 + -tmax + 1) *
32 < 0?-((-32 * c0 + -tmax + 1) / 32) : ((32 < 0?(-32 * c0 + -tmax + 1) + -32 - 1) / -32 : (32 * c0 + -tmax + 1 + 32
- 1) / 32))))?((c0 * 2 < 0?-(-c0 / 2) : ((2 < 0?(-c0 + -2 - 1) / -2 : (c0 + 2 - 1) / 2)))) : (((32 * c0 + -tmax + 1) *
32 < 0?-((-32 * c0 + -tmax + 1) / 32) : ((32 < 0?(-32 * c0 + -tmax + 1) + -32 - 1) / -32 : (32 * c0 + -tmax + 1 + 32
- 1) / 32))))); c1 <= (((((((ny + tmax + -2) * 32 < 0?((32 < 0?-((-ny + tmax + -2) + 32 + 1) / 32) : -((-ny + tmax +
-2) + 32 - 1) / 32))) : (ny + tmax + -2) / 32)) < (((32 * c0 + ny + 30) * 64 < 0?((64 < 0?-((-32 * c0 + ny + 30) + 64
+ 1) / 64) : -((-32 * c0 + ny + 30) + 64 - 1) / 64))) : (32 * c0 + ny + 30) / 64))?(ny + tmax + -2) * 32 < 0?((32 <
0?-((-ny + tmax + -2) + 32 + 1) / 32) : -((-ny + tmax + -2) + 32 - 1) / 32))) : (ny + tmax + -2) / 32)) : (((32 * c0 +
ny + 30) * 64 < 0?((64 < 0?-((-32 * c0 + ny + 30) + 64 + 1) / 64) : -((-32 * c0 + ny + 30) + 64 - 1) / 64))) : (32 *
c0 + ny + 30) / 64)) < c0?((((ny + tmax + -2) * 32 < 0?((32 < 0?-((-ny + tmax + -2) + 32 + 1) / 32) : -((-ny + tmax
+ -2) + 32 - 1) / 32))) : (ny + tmax + -2) / 32)) < (((32 * c0 + ny + 30) * 64 < 0?((64 < 0?-((-32 * c0 + ny + 30) + 64
+ 1) / 64) : -((-32 * c0 + ny + 30) + 64 - 1) / 64))) : (32 * c0 + ny + 30) / 64))?(ny + tmax + -2) * 32 < 0?((32 <
0?-((-ny + tmax + -2) + 32 + 1) / 32) : -((-ny + tmax + -2) + 32 - 1) / 32))) : (ny + tmax + -2) / 32)) : (((32 * c0 +
ny + 30) * 64 < 0?((64 < 0?-((-32 * c0 + ny + 30) + 64 + 1) / 64) : -((-32 * c0 + ny + 30) + 64 - 1) / 64))) : (32 *
c0 + ny + 30) / 64))))) : c0)); ++c1) {
    for (c2 = c0 + -c1; c2 <= (((((tmax + nx + -2) * 32 < 0?((32 < 0?-((-tmax + nx + -2) + 32 + 1) / 32) : -((-tmax +
nx + -2) + 32 - 1) / 32))) : (tmax + nx + -2) / 32)) < (((32 * c0 + -32 * c1 + nx + 30) * 32 < 0?((32 < 0?-((-32 * c0 +
-32 * c1 + nx + 30) + 32 + 1) / 32) : -((-32 * c0 + -32 * c1 + nx + 30) + 32 - 1) / 32))) : (32 * c0 + -32 * c1 + nx +
30) / 32))?(tmax + nx + -2) * 32 < 0?((32 < 0?-((-tmax + nx + -2) + 32 + 1) / 32) : -((-tmax + nx + -2) + 32 - 1) /
32))) : (tmax + nx + -2) / 32)) : (((32 * c0 + -32 * c1 + nx + 30) * 32 < 0?((32 < 0?-((-32 * c0 + -32 * c1 + nx + 30)
+ 32 + 1) / 32) : -((-32 * c0 + -32 * c1 + nx + 30) + 32 - 1) / 32))) : (32 * c0 + -32 * c1 + nx + 30) / 32))))); ++c2)
    {
        if (c0 == 2 * c1 && c0 == 2 * c2) {
            for (c3 = 16 * c0; c3 <= ((tmax + -1 < 16 * c0 + 30?tmax + -1 : 16 * c0 + 30)); ++c3)
                if (c0 % 2 == 0)
                    (ey[0])[0] = (_edge_[c3]);
            ..... (200 more lines!)
```

Performance gain: 2-6× on modern multicore platforms

# RoadMap for Tutorial: Math for Irregular Codes

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## Concepts

- Polyhedral model review
- Sparse computations as union of dense computations
- Sparse Polyhedral Framework (SPF)
- Polyrec

## Hands On Tutorial Goals

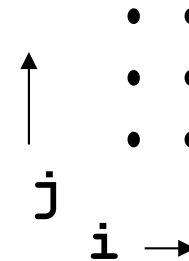
- Specify affine transformations in ISCC
- Handle irregular loop bounds in ISCC
- Demo of data dependence analysis for SPF
- Demo polyrec

# Representing Loops with Math (Matrices)

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## Original code

```
do i = 1,2
  do j = 1,3
    S1: A(i,j) = A(i-1,j+1)+1
  enddo
enddo
```



## Represent the iteration space

- As an intersection of inequalities
- The iteration space is the integer tuples within the intersection

**Bounds:**

$$1 \leq i \leq 2$$

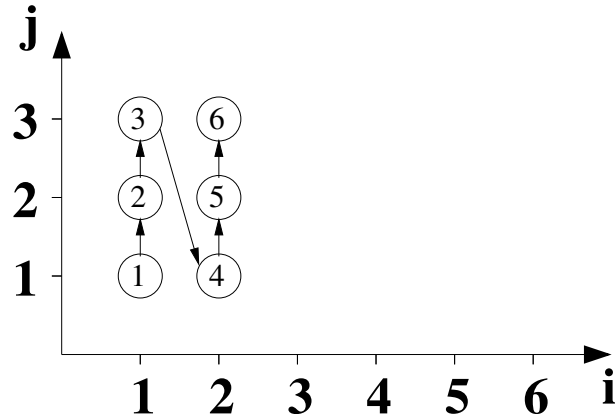
$$1 \leq j \leq 3$$

$$\begin{bmatrix} 1 & 0 \\ -1 & 0 \\ 0 & 1 \\ 0 & -1 \end{bmatrix} \begin{bmatrix} i \\ j \end{bmatrix} \geq \begin{bmatrix} 1 \\ -2 \\ 1 \\ -3 \end{bmatrix}$$

# Affine Transformations

## Interchange Transformation

The transformation matrix is the identity with a permutation of two rows.



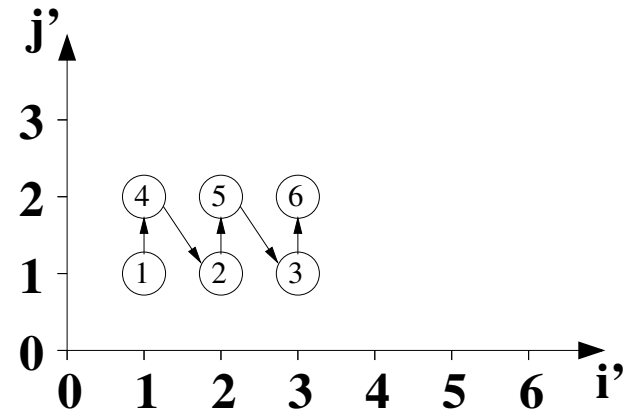
$$\begin{bmatrix} 1 & 0 \\ -1 & 0 \\ 0 & 1 \\ 0 & -1 \end{bmatrix} \begin{pmatrix} i \\ j \end{pmatrix} + \begin{pmatrix} -1 \\ 2 \\ -1 \\ 3 \end{pmatrix} \geq \vec{0}$$

(a) original polyhedron  
 $A\vec{x} + \vec{a} \geq \vec{0}$

$\Rightarrow$

$$\begin{pmatrix} i' \\ j' \end{pmatrix} = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} \begin{pmatrix} i \\ j \end{pmatrix}$$

(b) transformation function  
 $\vec{y} = T\vec{x}$



$$\begin{bmatrix} 0 & 1 \\ 0 & -1 \\ 1 & 0 \\ -1 & 0 \end{bmatrix} \begin{pmatrix} i' \\ j' \end{pmatrix} + \begin{pmatrix} -1 \\ 2 \\ -1 \\ 3 \end{pmatrix} \geq \vec{0}$$

(c) target polyhedron  
 $(AT^{-1})\vec{y} + \vec{a} \geq \vec{0}$

do  $i = 1, 2$   
do  $j = 1, 3$   
S( $i, j$ )

do  $i' = 1, 3$   
do  $j' = 1, 2$   
S( $i=j', j=i'$ )



# Goal: Learn How to Use ISCC

## ISCC

- Calculator for ISL (Integer Set Library)
- <http://compsys-tools.ens-lyon.fr/iscc/>
- Author: Sven Verdoolaege
- See Barvinok documentation online for a user manual

ISCC Online Demonstrator [Barvinok documentation](#) [ISL documentation](#) [Barvinok sources](#)

**Input** Reset Execute

ISCC script

```
I := [N] -> {S[i,j] : 0<=i<N and 0<=j<i};
codegen I;

T := {S[i,j] -> S[j,i]};
codegen (T*I);
```

**Output**

ISCC output

```
for (int c0 = 1; c0 < N; c0 += 1)
  for (int c1 = 0; c1 < c0; c1 += 1)
    S(c0, c1);
for (int c0 = 0; c0 < N - 1; c0 += 1)
  for (int c1 = c0 + 1; c1 < N; c1 += 1)
    S(c1, c0);
```

**Summary**

Command	iscc
Execution time	0.014 s
Version	barvinok-0.39 isl-0.16.1 polylib-5.22.5 ntl-6.2.1

Output format:

Indent output (experimental)

# Specifying Example Loop in ISCC

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## Original Loop in C

```
for (i=0; i<N; i++)
  for (j=0; j<i; j++)
    A[i][j] = exp(i+j);
```

## Create a macro for statement in C

```
#define S(i,j) A[i][j] = exp((i)+(j))
for (i=0; i<N; i++)
  for (j=0; j<i; j++)
    S(i,j);
```

## Iterations in loops described as a Set in ISCC

```
I := [N] -> {S[i,j] : 0<=i<N and 0<=j<i};
```

# ISCC: Loop Interchange Transformation

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## Generate the loop bounds for the Set I

**// Input**

```
I := [N] -> {S[i,j] : 0<=i<N and 0<=j<i};  
codegen I;
```

**// Output**

```
for (int c0 = 1; c0 < N; c0 += 1)  
  for (int c1 = 0; c1 < c0; c1 += 1)  
    S(c0, c1);
```

## Generate after applying *Loop Interchange*

**// Input: Transformation function**

```
T := {S[i,j] -> [j,i]};  
codegen (T*I);
```

**// Output**

```
for (int c0 = 0; c0 < N - 1; c0 += 1)  
  for (int c1 = c0 + 1; c1 < N; c1 += 1)  
    S(c1, c0);
```

# ISCC Determines Old Iterators as Function of New Iterators

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## Original Loop Nests and Transformation

**// Input**

```
Domain := [N] -> {S1[0, j, 0] : 1 <= j < N;  
                  S2[1, k, 0] : 0 <= k < N-1; };
```

```
T_fusion := {S1[0, j, 0] -> [0, j, 0];  
            S2[1, k, 0] -> [0, k+1, 1] };
```

```
codegen (T_fusion*Domain);
```

## Resulting code with old iterators as function of new iterators

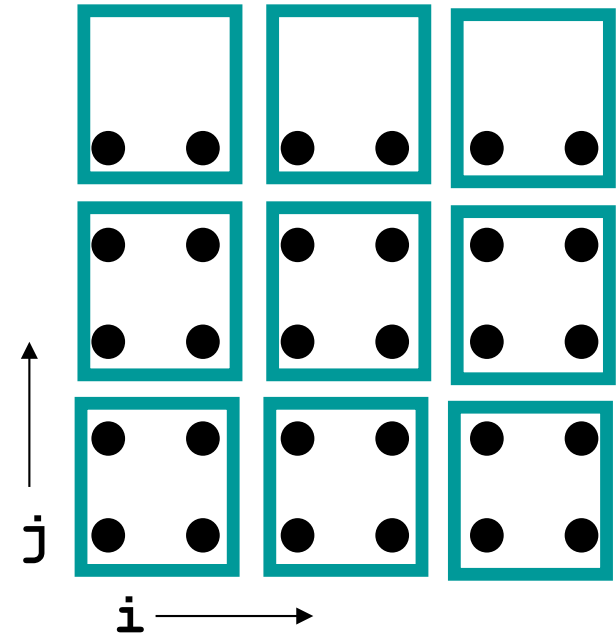
**// Output**

```
for (int c1 = 1; c1 < N; c1 += 1) {  
    S1(0, c1, 0);  
    S2(1, c1 - 1, 0);  
}
```

# Tiling

## A loop transformation that ...

- groups iteration points into tiles that are executed atomically
- can improve spatial and temporal data locality
- can expose larger granularities of parallelism



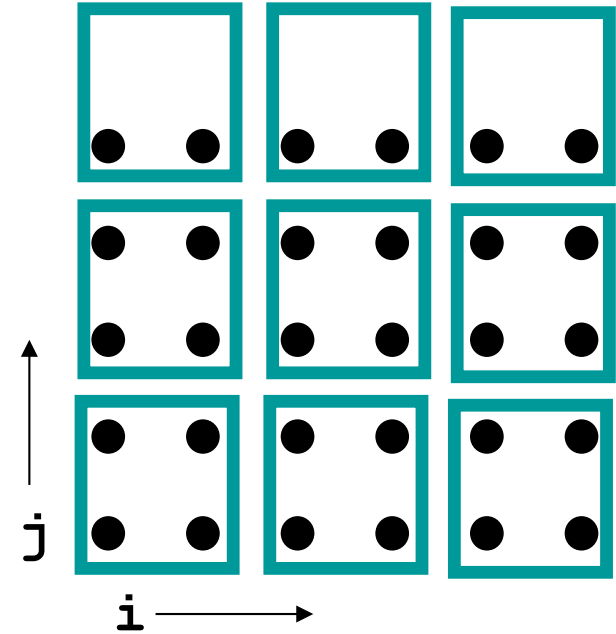
```
do ii = 1, 6, by 2
  do jj = 1, 5, by 2
    do i = ii, ii+2-1
      do j = jj, min(jj+2-1, 5)
        A(i, j) = ...
```

# Specifying Tiling

## Rectangular tiling

– tile size vector  $(s_1, s_2, \dots, s_d)$

– tile offset,  $(o_1, o_2, \dots, o_d)$



## Possible Transformation Mappings

– creating a tile space

$$\{[i, j] \rightarrow [ti, tj, i, j] \quad | \quad ti = \text{floor}((i - o_1)/s_1) \\ \wedge tj = \text{floor}((j - o_2)/s_2)\}$$

– keeping tile iterators in original iteration space

$$\{[i, j] \rightarrow [ii, jj, i, j] \quad | \quad ii = s_1 \text{floor}((i - o_1)/s_1) + o_1 \\ \wedge jj = s_2 \text{floor}((j - o_2)/s_2) + o_2\}$$

## Using ISCC to do code generation for tiling

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**Iteration space:**  $S := \{ s[i, j] : 1 \leq i \leq 6 \ \&\& \ 1 \leq j \leq 5 \};$

### Tiling specification

```
T := {s[i, j] -> [ti, tj, i, j] : ti = (i-1) / 2 && tj = (j-1) / 2};  
codegen (T*S); // doesn't work in iscc
```

### Getting rid of integer division

$ti = (i-1) / 2$  **becomes**

$0 \leq ri < 2 \ \&\& \ (i-1) = 2 * ti + ri$

$tj = (j-1) / 2$  **becomes**

$0 \leq rj < 2 \ \&\& \ (j-1) = 2 * tj + rj$

```
T := {s[i, j] -> [ti, tj, i, j] : exists ri, rj:
```

```
0 <= ri < 2 && i-1 = ti*2 + ri
```

```
&& 0 <= rj < 2 && j-1 = tj*2 + rj}; // works!!
```

# Polyhedral Model

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## Some History

- [Banerjee90] Uptal Banerjee, “Unimodular transformations of double loops,” In Advances in Languages and Compilers for Parallel Computing, 1990.
- [Wolf & Lam 91] Wolf and Lam, “A Data Locality Optimizing Algorithm,” In Programming Languages Design and Implementation, 1991.
- [Kelly and Pugh 95] Kelly and Pugh, “A unifying framework for iteration reordering transformations,” In IEEE First International Conference on Algorithms and Architectures for Parallel Processing (ICAPP)
- [Feautrier 96] Paul Feautrier, “Automatic Parallelization in the Polytope Model,” In The Data Parallel Programming Model.



# Polyhedral Model

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## Some key components

- Representing loops as sets
- Representing data dependences as dependence vectors
- Representing transformations as functions
- Applying transformations to generate transformed code